

# Minimum rank of a random graph over the binary field

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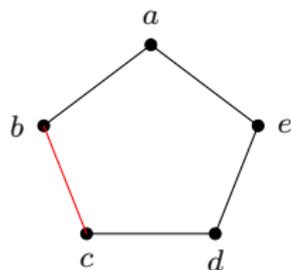
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# Definition (The minimum rank of a graph over a field)

A matrix  $M$  represents a graph  $G$  if

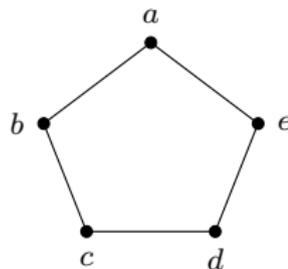
$$\begin{array}{c} a \\ b \\ c \\ d \\ e \end{array} \begin{array}{ccccc} & a & b & c & d & e \\ a & & 1 & 0 & 0 & 1 \\ b & 1 & & 1 & 0 & 0 \\ c & 0 & 1 & & 1 & 0 \\ d & 0 & 0 & 1 & & 1 \\ e & 1 & 0 & 0 & 1 & \end{array}$$



There are many matrices that represent a graph.  
Denote  $\text{mr}(\mathbb{F}, G)$ .

# Example ( $\text{mr}(\mathbb{F}_2, C_5) = 3$ )

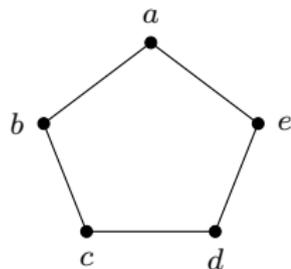
$$\begin{array}{c} a \\ b \\ c \\ d \\ e \end{array} \begin{pmatrix} a & b & c & d & e \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{pmatrix}$$



Thus,  $\text{mr}(\mathbb{F}_2, C_5) \leq 3$ .

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Thus,  $\text{mr}(\mathbb{F}_2, C_5) \geq 3$ .

an eigenvalue  $\lambda$   
of a matrix  $A$  which represents a graph  $G$

the (geometric) multiplicity  
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of an eigenvalue  $\lambda$   
 $= \text{nullity}(A - \lambda I)$

the maximum multiplicity of an eigenvalue  $\lambda$   
 $= \max \text{nullity}(A - \lambda I)$

$$\begin{aligned} & \max \text{ multiplicity of } \lambda \\ &= \max \text{ nullity}(A - \lambda I) \\ &= |V(G)| - \min \text{rank}(A - \lambda I) \\ &= |V(G)| - \text{mr}(G) \quad (\because A - \lambda I \text{ represents } G) \end{aligned}$$

Thus,

$$\text{mr}(G) = |V(G)| - \max \text{ multiplicity of } \lambda$$

## Some properties

- The minimum rank of  $G$  is at most 1 if and only if  $G$  can be expressed as the union of a clique and an independent set.
- A path  $P$  is the only graph of minimum rank  $|V(P)| - 1$ .
- For a cycle  $C$ ,  $\text{mr}(C) = |V(C)| - 2$ .
- If  $G'$  is an induced subgraph of  $G$ , then  $\text{mr}(G') \leq \text{mr}(G)$ .

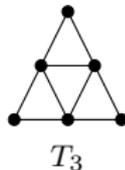
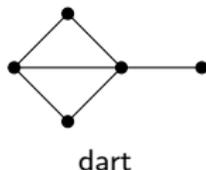
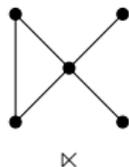
# Known results

Theorem(Barrett, van der Holst, and Loewy, 2004)

Let  $G$  be a graph. Then,  $\text{mr}(\mathbb{R}, G) \leq 2$  if and only if  $G$  is  $(P_4, \times, \text{dart}, P_3 \cup K_2, 3K_2, K_{3,3,3})$ -free.

Theorem(Hogben and van der Holst, 2006)

Let  $G$  be a 2-connected graph. Then,  $\text{mr}(\mathbb{R}, G) = n - 2$  if and only if  $G$  has no  $K_{4-}$ ,  $K_{2,3-}$ , or  $T_3$ -minor.



## Theorem(Ding and Kotlov, 2006)

If  $\mathbb{F}$  is a finite field, then for every  $k$ , the set of graphs of minimum rank at most  $k$  is characterized by finitely many forbidden induced subgraphs, each on at most  $\left(\frac{|\mathbb{F}|^k}{2} + 1\right)^2$  vertices.

## Remark

- $\text{mr}(\mathbb{F}_2, K_{3,3,3}) = 2$
- $\text{mr}(\mathbb{R}, K_{3,3,3}) = 3$

# Random graph

We consider the **Erdős-Rényi random graph**  $G(n, p)$ .

The vertex set of a random graph  $G(n, p)$  is  $\{1, 2, \dots, n\}$  and two vertices are adjacent with probability  $p$  independently at random.

Given a graph property  $\mathcal{P}$ , we say that  $G(n, p)$  possesses  $\mathcal{P}$  *asymptotically almost surely*, or a.a.s. for brevity, if the probability that  $G(n, p)$  possesses  $\mathcal{P}$  converges to 1 as  $n$  goes to infinity.

The minimum rank of a random graph over a field.

	$\mathbb{R}^\dagger$	$\mathbb{F}_2^\ddagger$
$G(n, 1/2)$	$0.147n < \text{mr} < 0.5n$	$n - \sqrt{2n} \leq \text{mr}$
$G(n, p)$	$cn < \text{mr} < dn$	

$\dagger$  Hall, Hogben, Martin, and Shader, 2010

$\ddagger$  Friedland and Loewy, 2010

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Let  $p(n)$  be a function s.t.  $0 < p(n) \leq \frac{1}{2}$  and  $np(n)$  is increasing. We prove that the minimum rank of  $G(n, 1/2)$  and  $G(n, p(n))$  over the binary field is at least  $n - o(n)$  a.a.s. We have two different proofs.

## Theorem (using the 1st method)

- $\text{mr}(\mathbb{F}_2, G(n, 1/2)) \geq n - \sqrt{2n} - 1.01$  a.a.s.
- $\text{mr}(\mathbb{F}_2, G(n, p(n))) \geq n - 1.483\sqrt{n/p(n)}$  a.a.s.  $(\sqrt{2 \ln 3})$

## Theorem (using the 2st method)

- $\text{mr}(\mathbb{F}_2, G(n, 1/2)) \geq n - 1.415\sqrt{n}$  a.a.s.
- $\text{mr}(\mathbb{F}_2, G(n, p(n))) \geq n - 1.178\sqrt{n/p(n)}$  a.a.s.  $(\sqrt{2 \ln 2})$

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## Theorem (J., C.Lee, P.Loh, S.Oum, 2013+)

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	$\mathbb{R}$	$\mathbb{F}_2$
$G(n, 1/2)$	$0.147n < \text{mr} < 0.5n$	$n - \sqrt{2n} \leq \text{mr}$
$G(n, p)$	$cn < \text{mr} < dn$ ( $p$ fixed)	$n - 1.178\sqrt{n/p(n)} \leq \text{mr}$

- A nontrivial **upper bound** of the minimum rank of a random graph over the binary field is an open question.
- The minimum rank of a random graph over **the other fields** is unknown.
- The minimum rank of a random graph  $G(n, p)$  is unknown.
- Is the minimum rank problem NP-complete??

Thank you.



Let  $p(n)$  be a function s.t.  $0 < p(n) \leq \frac{1}{2}$  and  $np(n)$  is increasing. We prove that the minimum rank of  $G(n, 1/2)$  and  $G(n, p(n))$  over the binary field is at least  $n - o(n)$  a.a.s. We have two different proofs.

## Theorem (using the 1st method)

- $\text{mr}(\mathbb{F}_2, G(n, 1/2)) \geq n - \sqrt{2n} - 1.01$  a.a.s. (Proof)
- $\text{mr}(\mathbb{F}_2, G(n, p(n))) \geq n - 1.483\sqrt{n/p(n)}$  a.a.s.

## Theorem (using the 2st method)

- $\text{mr}(\mathbb{F}_2, G(n, 1/2)) \geq n - 1.415\sqrt{n}$  a.a.s.
- $\text{mr}(\mathbb{F}_2, G(n, p(n))) \geq n - 1.178\sqrt{n/p(n)}$  a.a.s.

## Theorem

Let  $\mathbb{F}_2$  be the binary field and  $G(n, \frac{1}{2})$  be a random graph. Then,

$$\text{mr} \left( \mathbb{F}_2, G \left( n, \frac{1}{2} \right) \right) \geq n - \sqrt{2n} - 1.01$$

asymptotically almost surely.

## Sketch of the proof.

$G = G(n, 1/2)$

$\mathcal{G}_n$  : a set of all graphs with a vertex set  $\{1, 2, \dots, n\}$

$S_n(\mathbb{F}_2)$  : a set of all  $n \times n$  symmetric matrices over the binary field

There can be many different matrices representing the same graph. If one of them has rank less than  $r$ , then the minimum rank of this graph is less than  $r$ . Thus,

$$\sum_{\substack{\text{mr}(\mathbb{F}_2, H) < r \\ H \in \mathcal{G}_2}} \mathbb{P}[G = H] \leq \sum_{\substack{\text{rank}(N) < r \\ N \in \mathcal{M}}} \mathbb{P}[G = G(N)].$$

Let  $M$  be an  $n \times n$  random symmetric matrix s.t. every entry on or above the main diagonal of  $M$  is 1 with  $1/2$ . For  $N \in S_n(\mathbb{F}_2)$ , we have

$$\mathbb{P}[G = G(N)] = 2^n \mathbb{P}[M = N]$$

because the diagonal entries are decided with probability  $1/2$  independently at random.

Therefore, we have

$$\begin{aligned}\mathbb{P}[\text{mr}(\mathbb{F}_2, G) < n - L] &= \sum_{\substack{\text{mr}(\mathbb{F}_2, H) < n - L \\ H \in \mathcal{G}}} \mathbb{P}[G = H] \\ &\leq \sum_{\substack{\text{rank}(N) < n - L \\ N \in \mathcal{M}}} \mathbb{P}[G = G(N)] \\ &= 2^n \sum_{\substack{\text{rank}(N) < n - L \\ N \in \mathcal{M}}} \mathbb{P}[M = N] \\ &= 2^n \mathbb{P}[\text{rank}(M) < n - L] \\ &= 2^n \mathbb{P}[\text{nullity}(M) > L].\end{aligned}$$

It is enough to show that  $\mathbb{P}[\text{nullity}(M) > \sqrt{2n} + 1.0]$  is  $o(1/2^n)$ .  
So, we focus on  $\mathbb{P}[\text{nullity}(M) = L]$ .

## Lemma

Let  $M_i$  be an  $i \times i$  random symmetric matrix such that every entry in the upper triangle and diagonal of  $M_i$  is 1 with probability  $\frac{1}{2}$  independently at random. And let  $P_{i,k}$  be the probability that  $M_i$  has nullity  $k$ . Then,  $P_{1,0} = P_{1,1} = P_{2,0} = \frac{1}{2}$ ,  $P_{2,1} = \frac{3}{8}$ ,  $P_{2,2} = \frac{1}{8}$ ,  $P_{i,-1} = 0$  for all  $i$ ,  $P_{i,k} = 0$  for all  $i < k$ , and

$$P_{i,k} = \frac{1}{2}P_{i-1,k} + \frac{1}{2^i}P_{i-1,k-1} + \frac{1}{2}\left(1 - \frac{1}{2^{i-1}}\right)P_{i-2,k}$$

for  $i \geq 3$ ,  $k \geq 0$ .